Technological Inefficiency Indexes: A Binary Taxonomy and Generic Theorems

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Paper Prepared for the IARIW-UNSW Conference on Productivity: Measurement, Drivers and Trends

Sydney, Australia, November 26-27, 2013

Session 6B: Productivity and Efficiency Analysis

Time: Wednesday, November 27, 12:00-12:30
Draft

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October 2013

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PDF version of generic’31.10.13.tex produced on 3 November 2013.
I. Introductory Remarks.

Over the years, a large number of indexes of technological inefficiency (or, equivalently, technological efficiency) have been specified, and a spate of papers has examined the properties, or axioms, satisfied by these indexes. Russell and Schworm [2011] carried out a systematic analysis, axiom by axiom and specification by specification. Their theorems suggest, however, a more synthetic structure and the possibility of generic results on classes of indexes and their properties. The purpose of the present paper is to present such results. In particular, we consider a broad class of indexes containing almost all known indexes and a partition of this class into two subsets, which we term “slacks-based indexes” and “path-based indexes”. Slacks-based indexes are expressed in terms of additive or multiplicative slacks for all inputs and outputs, and particular indexes are generated by specifying the form of aggregation over the coordinate-wise slacks. Path-based indexes are expressed in terms of a common contraction/expansion factor, and particular indexes are generated by specifying the form of the path to the boundary of the technology.

Owing to an impossibility result of Russell and Schworm [2011], we know that the set of all inefficiency indexes can be partitioned into three subsets: those that satisfy continuity (in quantities and technologies) and violate indication (equal to some specified value if and only if the quantity vector is efficient), those that satisfy indication and violate continuity, and those that satisfy neither indication nor continuity. Abstracting from the indexes satisfying neither axiom, we prove two generic theorems showing the equivalence of these two taxonomies—i.e., showing that slacks-based indexes satisfy indication and hence violate continuity and path-based indexes satisfy continuity and hence violate indication. In our concluding remarks, we discuss briefly the few indexes that do not belong to either of these two sets.

II. Technology Sets and Efficiency.

The \((\text{input, output})\) production vector \(\langle x, y \rangle \in \mathbb{R}^{n+m}_+\) is constrained to lie in a technology set \(T \subset \mathbb{R}^{n+m}_+\). The output possibility set for input \(x\) is \(P(x) = \{y \in \mathbb{R}^m_+ \mid \langle x, y \rangle \in T\}\). Denote the origin of \(\mathbb{R}^{n+m}_+\) by \(0^{[n+m]} = (0^{[n]}, 0^{[m]})\) and the frontier of \(T\) by \(\partial T\).

We consider the collection of non-empty, closed technology sets, \(T\), that satisfy the following conditions:\footnote{Vector notation: \(\bar{x} \geq x\) if \(\bar{x}_i \geq x_i\) for all \(i\); \(\bar{x} > x\) if \(\bar{x}_i \geq x_i\) for all \(i\) and \(\bar{x} \neq x\); and \(\bar{x} \gg x\) if \(\bar{x}_i > x_i\) for all \(i\).}

\[
\begin{align*}
\text{(i) } & \langle x, y \rangle \in T \text{ and } \langle \bar{x}, -\bar{y} \rangle \geq \langle x, -y \rangle \text{ implies } \langle \bar{x}, \bar{y} \rangle \in T \text{ (free disposability of inputs and outputs),} \\
\text{(ii) } & y > 0^{[m]} \implies \langle 0^{[n]}, y \rangle \notin T \text{ (no free lunch), and}
\end{align*}
\]
(iii) $P(x)$ is non-empty and bounded for all $x \in \mathbb{R}_+^n$.

Note that, owing to condition (i), $\langle x, y \rangle \in \partial T$ only if $\langle x, -y \rangle \gg \langle \bar{x}, -\bar{y} \rangle$ implies $\langle \bar{x}, \bar{y} \rangle \notin T$.

Many inefficiency indexes were originally defined on the particular subset of $\mathcal{T}$ generated by mathematical programming methods of constructing technology sets on a finite set of data points. This method, commonly referred to as Data Envelopment Analysis (DEA), generates convex polyhedral technologies (i.e., intersections of finite numbers of half spaces). A most all of these indexes, however, can be applied to the more general class of technologies $\mathcal{T}$.

A production vector $\langle x, y \rangle \in \mathcal{T}$ is technologically efficient (in the sense of Koopmans [1951]) if $\langle x, -y \rangle > \langle \bar{x}, -\bar{y} \rangle$ implies $\langle \bar{x}, \bar{y} \rangle \notin \mathcal{T}$. Denote the set of efficient production vectors in $\mathcal{T}$ by $\text{Eff}(\mathcal{T})$.

III. Inefficiency Indexes: A Binary Taxonomy.

1. Technological efficiency indexes.

Intuitively, a technological inefficiency index measures the “distance” from the production vector to a “reference point” on the frontier of $\mathcal{T}$. Formally, we define an inefficiency index as a mapping, $I : \Xi \to \mathbb{R}^n$, with image $I(x, y, T)$, where

$$\Xi = \{ \langle x, y, T \rangle \in \mathbb{R}_+^{n+m} \times \mathcal{T} \mid \langle x, y \rangle \in \mathcal{T} \}$$

and $\mathbb{R}^n$, the effective range of $I$, is a closed subset of $\mathbb{R}_+^n$.

2. Slacks-based indexes.

Slacks-based indexes can be formulated in terms of either additive or multiplicative slacks: viz., the additive or multiplicative difference between a given production vector $\langle x, y \rangle$ and some reference vector on the frontier determined by the optimization problem in the definition of a specific efficiency index. We denote the additive slacks by $\langle s, t \rangle \in \mathbb{R}_+^n \times \mathbb{R}_+^m$. Proportional input and output slacks are then given respectively by $\alpha_i = (x_i - s_i)/x_i \in (0, 1]$ for all $i$ and $\beta_j = y_j/(y_j + t_j) \in (0, 1]$ for all $j$. To facilitate easy comparison, we formulate

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2 See Charnes, Cooper, Lewin, and Seiford, [1985]
3 It is a straightforward matter of renormalization to convert an efficiency measure (typically mapping into the (0,1] interval) into an inefficiency measure.
4 Note that, to focus on the salient issues at hand, we restrict the domain of $I$ to positive values of input and output quantities, thus avoiding some distracting boundary issues (see Levkoff, Russell, and Schworm [2012] for an analysis of boundary problems).
5 Of course, proportional slacks are conversely converted to additive slacks by $s_i = (1 - \alpha_i)x_i$ for all $i$ and $t_j = (1 - \beta_j)y_j$ for all $j$. 

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all slacks-based indexes in terms of additive slacks (although some were originally formulated in terms of multiplicative slacks).

A slacks-based inefficiency index is defined by

\[ I^s_b(x, y, T) = \max_{(s, t) \in \mathbb{R}_{n+m}^+} \{ \psi(s, t, x, y) \mid \langle x - s, y + t \rangle \in T \}, \]  

(3.2)

where the function \( \psi \) is independent of units of measurement and increasing in the slack variables \( (s, t) \) and satisfies \( \psi(0^{[n+m]}, x, y) = 0 \) for all \( \langle x, y \rangle \in T \). Members of this family of inefficiency indexes are generated by specifying the form of the function \( \psi \).

The following indexes—converted to inefficiency indexes and/or normalized to satisfy \( \psi(0^{[n+m]}, x, y) = 0 \) where necessary—satisfy the definition (3.2) with the indicated specifications of \( \psi \).\(^6\)

- **Färe-Grosskopf-Lovell Index** (Färe, Grosskopf, and Lovell [1985]):

\[ \psi(s, t, x, y) = \left[ \frac{1}{n + m} \left( \sum_i \frac{x_i - s_i}{x_i} + \sum_j \frac{y_j}{y_j + t_j} \right) \right]^{-1} - 1. \]  

(3.3)

- **Quantity-Weighted Additive Index** (Charnes, Cooper, Golany, Seiford, and Stutz [1985]):\(^7\)

\[ \psi(s, t, x, y) = \sum_i s_i \frac{x_i}{x_i} + \sum_j t_j \frac{y_j}{y_j}. \]  

(3.4)

- **Weighted Additive Index** (Cooper and Pastor [1995]):

\[ \psi(s, t, x, y) = \sum_i u_i s_i + \sum_j v_j t_j \]  

(3.5)

where \( u \in \mathbb{R}_n^+ \) and \( v \in \mathbb{R}_m^+ \) are pre-specified weights.

- **Pastor-Ruiz-Sirvent Index** (Pastor, Ruiz, and Sirvent [1999]):

\[ \psi(s, t, x, y) = \frac{1}{n} \sum_i (x_i - s_i)/x_i - \frac{1}{m} \sum_j (y_j + t_j)/y_j - 1. \]  

(3.6)

\(^6\) See Russell and Schworm [2011] for an in-depth comparison of these indexes and those that follow below.

\(^7\) Charnes, Cooper, Golany, Seiford, and Stutz also specified the “Additive DEA Model” (eschewing the weights), but that index is not independent of units of measurement and hence has been superceded by the Quantity-Weighted Additive Index.
• **Measure of Efficiency Proportions** (Banker and Cooper [1994]): equivalent to the Färe-Grosskopf-Lovell index.

• **Measure of Inefficiency Proportions** (Cooper, Park, and Pastor [1999]): equivalent to the Quantity-Weighted Additive Index.

• **Slacks-Based Measure of Efficiency** (Tone [2001]): equivalent to the Pastor-Ruiz-Sirvent Index.

• **Directional Slacks-Based Measure** (Fukuyama and Weber [2009]): equivalent to the Weighted Additive Index.

3. **Path-based indexes.**

Path-based indexes are defined by

\[ I^{pb}(x, y, T) = \max \{ \lambda \in \Lambda \mid \omega(x, y, \lambda) \in T \}, \]  

(3.7)

where \( \Lambda \) is a closed (unit free) subset of \( \mathbb{R}_+ \), with \( \theta = \min \Lambda \), and the contraction/expansion mapping \( \omega \) from \( T \times \Lambda \) into \( T \) is continuous and satisfies

(i) \( \omega(x, y, \lambda) = \omega(x', y', \lambda') \) and \( \lambda > \lambda' \) imply \( \langle x, -y \rangle \gg \langle x', -y' \rangle \)

and

(ii) \( \lim_{\lambda \to \infty} \omega_j(x, y, \lambda) = \infty \) for some output \( j \).

For a given \( \langle \bar{x}, \bar{y} \rangle \), \( \{ \omega(\bar{x}, \bar{y}, \lambda) \mid \lambda \in \Lambda \} \) traces out the contraction/expansion path from \( \langle \bar{x}, \bar{y} \rangle \) to the boundary of \( T \), and \( \omega(\bar{x}, \bar{y}, I^{pb}(\bar{x}, \bar{y}, T)) \in \partial(T) \) is the reference vector for \( \langle \bar{x}, \bar{y} \rangle \). Restriction (i) says that, along level curves in \( \mathbb{R}^2_+ \) where \( \omega(x, y, \lambda) \) is a constant vector, an increase in \( \lambda \) requires an increase in \( x_i \) for all \( i \) and a decrease in \( y_j \) for all \( j \); in other words, increasing \( \lambda \) contracts \( x \) and expands \( y \). Restriction (ii) says that, as the contraction/expansion factor is increased indefinitely, at least one output quantity must expand indefinitely.

The following indexes satisfy the definition (3.7) with the indicated specifications of \( \Lambda \) and \( \omega \).

• **The Hyperbolic Index** (Färe, Grosskopf, and Lovell [1985]):

\[ \Lambda = [1, +\infty) \quad \text{and} \quad \omega(x, y, \lambda) = \langle x/\lambda, \lambda y \rangle. \]  

(3.8)
• *The Directional-Distance Index* (Luenberger [1992] and Chung, Färe, and Grosskopf [1997]8):

\[ \Lambda = [0, \infty) \text{ and } \omega(x, y, \lambda) = \langle x - \lambda g_x, y + \lambda g_y \rangle, \text{ where } g = \langle g_x, g_y \rangle \in \mathbb{R}_+^{n+m}. \]

\[ (3.9) \]

• *The Briec [1997] Index*9:

\[ \Lambda = [0, 1] \text{ and } \omega(x, y, \lambda) = \langle (1 - \lambda)x, (1 + \lambda)y \rangle. \]

\[ (3.10) \]

IV. Generic Theorems.

Among the axioms, or properties, of efficiency and inefficiency indexes that have been posited in the literature, two are salient for the above taxonomy:

*Indication of efficiency* (I): For all \( \langle x, y, T \rangle \in \Xi \), there exists a \( \theta \in R(I) \) such that \( I(x, y, T) = \theta \) if and only if \( \langle x, y \rangle \in \text{Eff}(T) \).

*Joint continuity* (C): \( I \) is continuous in \( \langle x, y, T \rangle \).

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The intrinsic appeal of these axioms is self-evident. Condition (I), introduced in the context of input-oriented indexes by Färe and Lovell [1978], requires that an inefficiency index distinguish between inefficient and efficient production vectors. In formulating condition (C), Russell [1990] argued (page 256) that continuity is a compelling property, “for it provides assurance that ‘small’ errors of measurement (of, e.g., input or output quantities) result only in ‘small’ errors of efficiency measurement.” If the technology is constructed (with error) from data on production vectors, the argument for continuity in the technology is perhaps even more compelling.

These two properties take on more importance when one considers the following result:

**Fact 1** (Russell and Schworm [2011, Theorem 1]): There does not exist an inefficiency index satisfying (I) and (C) on \( T \).

Thus, the set of all inefficiency indexes can be partitioned into three sets: (i) those that satisfy the indication property, (ii) those that satisfy continuity, and (iii) those that

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8 The directional distance index is adapted from the shortage function of Luenberger [1992] to the measurement of efficiency by Chung, Färe, and Grosskopf [1997]. Both restrict \( g \) only to the non-negative orthant, but restriction to the positive orthant improves the properties of the index (see Russell and Schworm [2011]).

9 Briec [1997] derives this index from the directional-distance function by using the definition of \( I_{DD} \) with the direction \( g = \langle x, y \rangle \).

10 As in Russell [1990] and Russell and Schworm [2011], we adopt the topology of closed convergence on \( T \).
satisfy neither. For reasons we discuss briefly in the Section 5, we have little interest in those satisfying neither. Instead, we are interested in characterizing a binary partition of indexes—extant or yet to be formulated—into those satisfying indication and those satisfying continuity. To this end, we show that all indexes with the structure (3.2) satisfy the indication property but violate continuity, whereas all indexes with the structure (3.7) satisfy continuity but violate the indication property.

**Theorem 1:** $I^{sb}(x, y, T)$ violates (C) and satisfies (I).

**Proof:** Given Fact 1, it suffices to prove satisfaction of (I).

Suppose that $⟨x, y⟩ ∈ T$, with $⟨x, y⟩ \gg 0^{[n+m]}$, is not efficient so that there exists a production vector $⟨x', y'⟩ ∈ T$ satisfying $⟨x', y'⟩ \gg 0^{[n+m]}$ and $⟨x', -y'⟩ < ⟨x, -y⟩$. If

$$⟨s', t'⟩ = \underset{(s, t) \in R^{n+m}}{\text{argmax}} \{ψ(x', y', s, t) | ⟨x' - s, y' + t⟩ ∈ T\},$$

$⟨x' - s', y' + t'⟩$ is a boundary point of $T$. As $⟨x' - s', -(y' + t')⟩ \leq ⟨x', -y'⟩ < ⟨x, -y⟩$, there exists an $⟨s^o, t^o⟩ ∈ R^{n+m}$ such that $⟨s^o, t^o⟩ > ⟨s', t'⟩$ and $⟨x - s^o, y + t^o⟩ = ⟨x' - s', y' + t'⟩ ∈ T$. As $ψ$ is increasing in $⟨s, t⟩$, it must be that $I^{sb}(x, y, T) > I^{sb}(x', y', T) ≥ 0$.

Next suppose that $I^{sb}(x, y, T) > 0$. Then $⟨s, t⟩ > 0^{[n+m]}$, so that there exists a point $⟨x - s, y + t⟩ =: ⟨x', y'⟩ ∈ T$ satisfying $⟨x', y'⟩ \gg 0^{[n+m]}$ and $⟨x', -y'⟩ < ⟨x, -y⟩$. Therefore, $(x, y)$ is inefficient. ■

**Theorem 2:** $I^{pb}(x, y, T)$ violates (I) and satisfies (C).

**Proof:** Given Fact 1, it suffices to prove satisfaction of (C).

Consider a sequence $\{x^{ν}, y^{ν}, T^{ν}\}$ converging to $\{x^o, y^o, T^o\}$ and let $ω^{ν} = ω(x^{ν}, y^{ν}, λ^{ν})$ and $ω^o = ω(x^o, y^o, λ^o)$, where $λ^{ν} = I^{pb}(x^{ν}, y^{ν}, T^{ν})$, and $λ^o = I^{pb}(x^o, y^o, T^o)$. The strategy is to show that, for arbitrary $ε$, there exists a $ν'$ such that

$$λ^o - ε < λ^{ν} < λ^o + ε \quad ∀ ν > ν'.$$  

Given $x^{ν} → x^o$ and $T^{ν} → T^o$, it follows that $P(x^{ν}) → P(x^o)$. By assumption, $P(x^0)$ is bounded; consequently, there exists a $ν'$ such that $w^{ν}$ is bounded for all $ν > ν'$. Invoking (ii), it follows that $λ^{ν}$ is bounded for all $ν > ν'$.

To prove the second inequality in (4.2), suppose the contrary, so that $λ^{ν} ≥ λ^o + ε$ for infinitely many $ν$ and some $ε > 0$. As $λ^{ν}$ is bounded for sufficiently large $ν$, this sequence has a convergent subsequence: $λ^{ν_κ} → ̄λ > λ^o$. Since $w^{ν_κ} ∈ T^{ν_κ}k$ for all $ν_κ$ and $T^{ν} → T^o$, we have $ω(x^o, y^o, ̄λ) ∈ T^o$. Along with the definition of $I^{pb}$ in (3.7), this implies that $̄λ ≤ λ^o$, a contradiction.
To prove the first inequality in (4.2), suppose the contrary so that \( \lambda^{\nu} \leq \lambda^o - \epsilon \) for infinitely many \( \nu \) and some \( \epsilon > 0 \). As \( \lambda^{\nu} \) is bounded for sufficiently large \( \nu \), this sequence has a convergent subsequence: \( \lambda^{\nu_k} \rightarrow \bar{\lambda} < \lambda^o \). As \( T^{\nu} \rightarrow T^o \) and \( \omega \) is continuous, it follows that \( w^{\nu_k} \), a boundary point in \( T^{\nu_k} \), converges to a boundary point, \( \bar{w} \). Using \( \bar{\lambda} < \lambda^o \) and property (i), this implies that \( \bar{w}_x \gg w^o_x \) and \( w^o_y \ll \bar{w}_y \). As \( \bar{w} \) and \( w^o \) are boundary points of \( T^o \), this violates the free disposability assumption. □

V. Concluding Remarks.

Axioms other than indication and continuity have been explored in the literature, but none seems as dispositive as these two. Invariance with respect to units of measurement has been built into the definitions of slacks-based and path-based indexes (3.2) and (3.7), and all indexes in the literature satisfy this fundamental condition.\(^{11}\)

Similarly, slacks-based and path-based indexes both satisfy (weak) monotonicity: \( \langle x, y \rangle \in T, \langle x', y' \rangle \in T, \) and \( \langle x', -y' \rangle \geq \langle x, -y \rangle \) imply \( I(x', y', T) \geq I(x, y, T) \). This is evident from the definitions of the two types of indexes. If \( \langle s^o, t^o \rangle \) solves the optimization problem in (3.2) at \( \langle x, y, T \rangle \) and \( \langle x', -y' \rangle \geq \langle x, -y \rangle \), then (by free disposability) \( \langle x' - s^o, y' + t^o \rangle \in T \); consequently \( \langle s^o, t^o \rangle \) is a feasible solution at \( \langle x', y', T \rangle \), implying that \( I^{sb}(x', y', T) \geq I^{sb}(x, y, T) \). Similarly, if \( \lambda^o \) solves the optimization problem in (3.7) at \( \langle x, y, T \rangle \) and \( \langle x', -y' \rangle \geq \langle x, -y \rangle \), then, using free disposability and condition (i), \( \omega(x', y', \lambda^o) \in T \), implying that \( I^{pb}(x', y', T) \geq I^{pb}(x, y, T) \).\(^{12}\)

Many efficiency studies, including those employing data envelopment analysis, restrict the collection of allowable technologies to convex polyhedral sets. It can be shown that, on this restricted class of technologies, slacks-based inefficiency indexes are continuous in input and output quantities.\(^{13}\) Nevertheless, Theorem 1 (and, a fortiori, Theorem 2) remains valid on this restricted class of technologies if the continuity condition is replaced by continuity in technologies. Thus, the indication-continuity trade-off between slacks-based and path-based technologies remains.

\(^{11}\) Other than the “Additive Model” of Charnes, Cooper, Golany, Seiford, and Stutz [1985], which has been superseded by the Weighted Additive Index formulated in the same paper for the explicit purpose of establishing unit invariance.

\(^{12}\) It is not possible to show that either class of indexes satisfies strict monotonicity: \( \langle x, y \rangle \in T, \langle x', y' \rangle \in T, \) and \( \langle x', -y' \rangle \geq \langle x, -y \rangle \) imply \( I(x', y', T) > I(x, y, T) \). One particular specification, the Weighted Additive Index, does satisfy this condition, but it has the disadvantage—admittedly not part of our axiomatic structure—of requiring the use of arbitrary weights to correct for the dependence of the “Additive Model” on unit changes.

\(^{13}\) The proof follows closely the continuity proof of selected slacks-based indexes in Russell and Schworm [2011, pp. 155–156].
Researchers have specified inefficiency indexes that do not belong to either the slacks-based or path-based family. The most interesting is the Weighted Holder Distance Function (Briec [1998]), which is based explicitly on a mathematical distance function. This index has the same properties as the path-based indexes but to our knowledge has yet to be applied in practice.\(^{14}\) Other indexes that do not fit into our taxonomy have distinctively inferior properties. For example, the multi-stage indexes (\textit{e.g.}, Coelli [1998]), sequentially combining path-based and slacks-based properties, appear to violate (weak) monotonicity (see the discussion of the Zieschang [1984] index in Russell and Schworm [2009]). Indexes that concomitantly combine slacks-based and path-based properties (\textit{e.g.}, the generalized hyperbolic measure of Färe, Grosskopf, and Lovell [1985]) appear to combine the worst features of both types of indexes, violating both continuity and identification.

Both types of indexes can be restricted to a subspace of \(\mathbb{R}^{n+m}_{+}\) (to generate, \textit{e.g.}, input-oriented indexes and output-oriented indexed). Subspace adaptations of the path-based efficiency indexes retain continuity, but subspace specializations of slacks-based indexes lose indication, since efficiency in a subspace does not imply efficiency in the full space of inputs and outputs (though one can, of course, relax the indication axiom to encompass only subspace efficiency).

To summarize, the axiomatic analysis of inefficiency measurement seems to boil down to the choice between the indication and continuity properties. The incompatibility of these two properties poses a fundamental quandary for the researcher. The trade-off between the slacks-based and the path-based indexes evident in Theorems 1 and 2 reflects the trade-off between these two properties. If the researcher values indication more than continuity, employ or specify a slacks-based index; if the researcher values continuity more than indication, employ or specify a path-based index.

We close by paying tribute to the formulations in two landmark papers that crystallize the fundamental dichotomy, and resultant specification quandary, of inefficiency indexes and axioms. The Farrell [1956] (input based) index is the model for path-based indexes (satisfying continuity but violating indication), and the Färe-Lovell [1978] (input based) index is the model for slacks-based indexes (satisfying indication but violating continuity). All subsequently specified (in)efficiency indexes are variations on these two themes.\(^{15}\)

\(^{14}\) The definition of path-based indexes (3.7) could be modified to encompass this index—after all, the Holder distance identifies the shortest path from a production vector to the boundary of the technology—but the redefinition would complicate the analysis (and would be a bit contrived).

\(^{15}\) In fact, the hyperbolic (in)efficiency index (3.8) is a straightforward generalization of the Farrell index (which takes a radial path to the boundary of the input-requirement set), and the Färe-Grosskopf-Lovell index (3.3) is a straightforward generalization of the Färe-Lovell index.
References


